

# CS 33

## Introduction to C Part 7

# The String Library

```
#include <string.h>

char *strcpy(char *dest, char *src);
    // copy src to dest, returns ptr to dest
char *strncpy(char *dest, char *src, int n);
    // copy at most n bytes from src to dest
int strlen(char *s);
    // returns the length of s (not counting the null)
int strcmp(char *s1, char *s2);
    // returns -1, 0, or 1 depending on whether s1 is
    // less than, the same as, or greater than s2
int strncmp(char *s1, char *s2, int n);
    // do the same, but for at most n bytes
```

---

# The String Library (more)

```
size_t strspn(const char *s, const char *accept);  
    // return length of initial portion of s  
    // consisting entirely of bytes from accept
```

```
size_t strcspn(const char *s, const char *reject);  
    // return length of initial portion of s  
    // consisting entirely of bytes not from  
    // reject
```

# Quiz 1

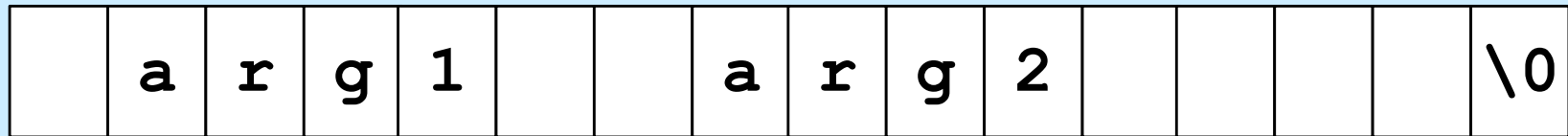
```
#include <stdio.h>
#include <string.h>

int main() {
    char s1[] = "Hello World!\n";
    char *s2;
    strcpy(s2, s1);
    printf("%s", s2);
    return 0;
}
```

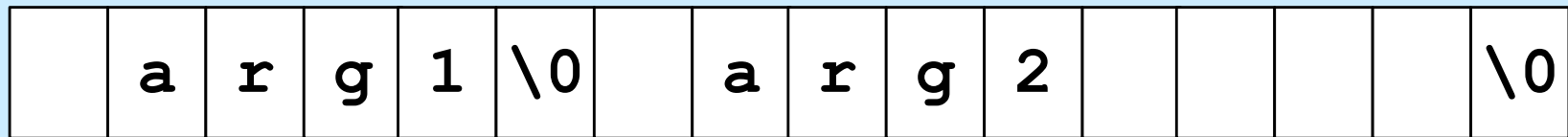
**This code:**

- a) has syntax problems**
- b) might seg fault**
- c) is a great example of well written C code**

# Parsing a String

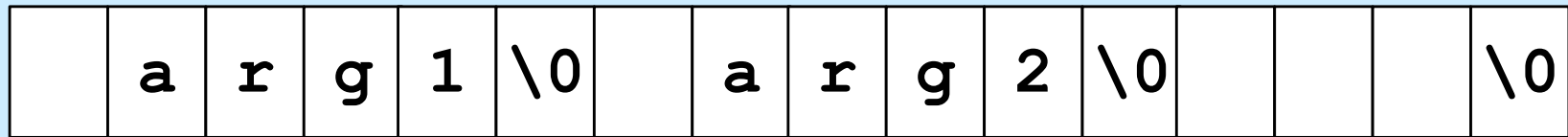


string



token

rem



token

rem

# Designing the Parse Function

- **It modifies the string being parsed**
  - puts nulls at the end of each token
- **Each call returns a pointer to the next token**
  - how does it know where it left off the last time?
    - » how is *rem* dealt with?

# Design of *strtok*

- **char** \**strtok*(**char** \**string*,  
                  **const char** \**sep*)
  - if *string* is non-NULL, *strtok* returns a pointer to the first token in *string* (and keeps track of where the next token would be)
  - if *string* is NULL, *strtok* returns a pointer to the token just after the one returned in the previous call, or NULL if there are no more tokens
  - tokens are separated by any non-empty combination of characters in *sep*

# Using *strtok*

```
int main() {
    char line[] = " arg0 arg1 arg2 arg3 ";
    char *str = line;
    char *token;
    while ((token = strtok(str, " \t\n")) != NULL) {
        printf("%s\n", token);
        str = NULL;
    }
    return 0;
}
```

## Output:

```
arg0
arg1
arg2
arg3
```



# *strtok* Code part 1

```
char *strtok(char *string, const char *sep) {
    static char *rem = NULL;
    if (string == NULL) {
        if (rem == NULL) return NULL;
        string = rem;
    }
    int len = strlen(string);
    int slen = strspn(string, sep);
    // initial separators
    if (slen == len) {
        // string is all separators
        rem = NULL;
        return NULL;
    }
}
```

## ***strtok* Code part 2**

```
string = &string[slen]; // skip over separators
len -= slen;
int tlen = strcspn(string, sep); // length of first token
if (tlen < len) {
    // token ends before end of string: terminate it with 0
    string[tlen] = '\0';
    rem = &string[tlen+1];
} else {
    // there's nothing after this token
    rem = NULL;
}
return string;
}
```

# Numeric Conversions

```
short a;
```

```
int b;
```

```
float c;
```

```
b = a;    /* always works */
```

```
a = b;    /* sometimes works */
```

```
c = b;    /* sort of works */
```

```
b = c;    /* sometimes works */
```

# Implicit Conversions (1)

```
float x, y=2.0;
```

```
int i=1, j=2;
```

```
x = i/j + y;
```

```
/* what's the value of x? */
```

# Implicit Conversions (2)

```
float x, y=2.0;
```

```
int i=1, j=2;
```

```
float a, b;
```

```
a = i;
```

```
b = j;
```

```
x = a/b + y;
```

```
/* now what's the value of x? */
```

# Explicit Conversions: Casts

```
float x, y=2.0;
```

```
int i=1, j=2;
```

```
x = (float)i / (float)j + y;
```

```
/* and now what's the value of x? */
```

# Purposes of Casts

- **Coercion**

```
int i, j;  
float a;  
a = (float) i / (float) j;
```

modify the  
value  
appropriately

- **Intimidation**

```
float x, y;  
// sizeof(float) == 4  
swap((int *) &x, (int *) &y);
```

it's ok as is  
(trust me!)

# Quiz 2

- Will this work?

```
double x, y; //sizeof(double) == 8
```

```
...
```

```
swap((int *) &x, (int *) &y);
```

- a) yes
- b) no



# Caveat Emptor

- **Casts tell the C compiler:**  
“Shut up, I know what I’m doing!”
- **Sometimes true**

```
float x, y;  
swap((int *) &x, (int *) &y);
```

- **Sometimes false**

```
double x, y;  
swap((int *) &x, (int *) &y);
```

# Nothing, and More ...

- ***void* means, literally, nothing:**

```
void NotMuch(void) {  
    printf("I return nothing\n");  
}
```

- **What does *void \** mean?**
  - it's a pointer to anything you feel like
    - » a generic pointer

# Rules

- **Use with other pointers**

```
int *x;
```

```
void *y;
```

```
x = y; /* legal */
```

```
y = x; /* legal */
```

- **Dereferencing**

```
void *z;
```

```
func(*z); /* illegal!*/
```

```
func(*(int *)z); /* legal */
```

# Swap, Revisited

```
void swap(int *i, int *j) {  
    int tmp;  
    tmp = *j; *j = *i; *i = tmp;  
}  
/* can we make this generic? */
```

# An Application: Generic Swap

```
void gswap (void *p1, void *p2,
           int size) {
    int i;
    for (i=0; i < size; i++) {
        char tmp;
        tmp = ((char *)p1)[i];
        ((char *)p1)[i] = ((char *)p2)[i];
        ((char *)p2)[i] = tmp;
    }
}
```

# Using Generic Swap

```
short a=1, b=2;  
gswap (&a, &b, sizeof(short));
```

```
int x=6, y=7;  
gswap (&x, &y, sizeof(int));
```

```
int A[] = {1, 2, 3}, B[] = {7, 8, 9};  
gswap (A, B, sizeof(A));
```

# Fun with Functions (1)

```
void ArrayDouble(int A[], int len) {  
    int i;  
    for (i=0; i<len; i++)  
        A[i] = 2*A[i];  
}
```

# Fun with Functions (2)

```
void ArrayBop(int A[],  
             int len,  
             int (*func)(int)) {  
    int i;  
    for (i=0; i<len; i++)  
        A[i] = (*func)(A[i]);  
}
```



# Fun with Functions (3)

```
int triple(int arg) {  
    return 3*arg;  
}
```

```
int main() {  
    int A[20];  
    ... /* initialize A */  
    ArrayBop(A, 20, triple);  
    return 0;  
}
```

# typedef

- **Allows one to create new names for existing types**

```
typedef int *IntP_t;
```

```
IntP_t x;
```

– means the same as

```
int *x;
```

# More typedefs

```
typedef struct complex {  
    float real;  
    float imag;  
} complex_t;
```

```
complex_t i, *ip;
```

# Not a Quiz

- What's A?

```
typedef double X_t[N];  
X_t A[M];
```

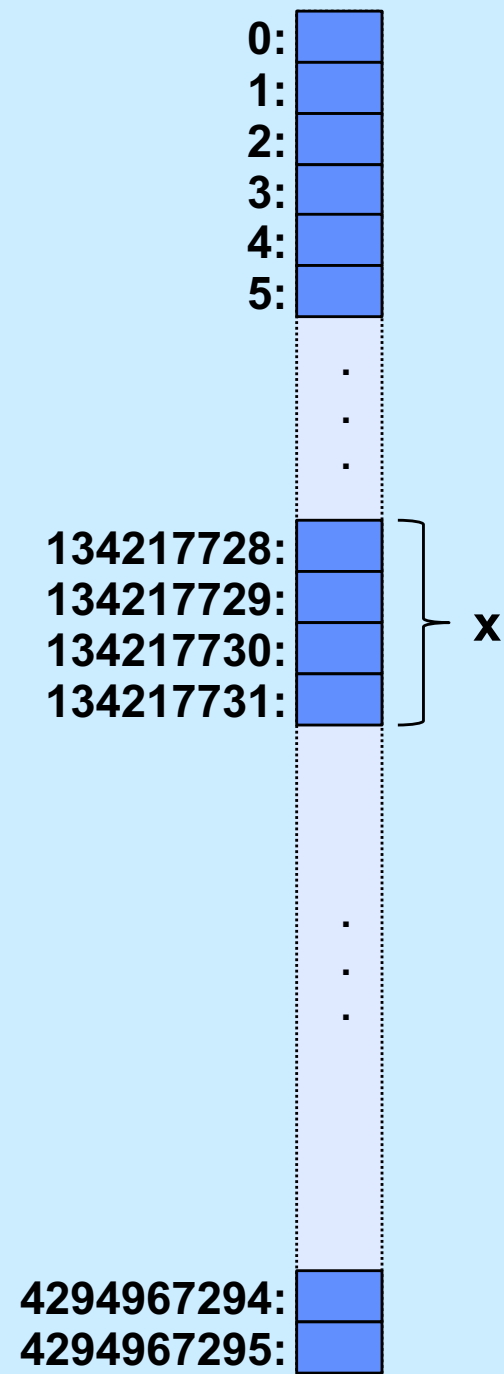
- a) an array of M doubles
- b) an MxN array of doubles
- c) an NxM array of doubles
- d) a syntax error

# CS 33

## Data Representation, Part 1

# Representing Data in Memory

- **x** is a 4-byte integer
  - how do the 32 bits represent its value?



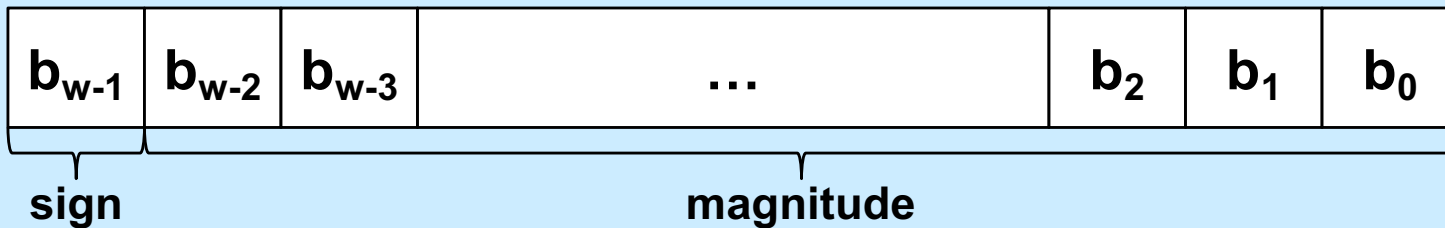
# Unsigned Integers



$$\text{value} = \sum_{i=0}^{w-1} b_i \cdot 2^i$$

# Signed Integers

- **Sign-magnitude**



$$\text{value} = (-1)^{b_{w-1}} \cdot \sum_{i=0}^{w-2} b_i \cdot 2^i$$

- two representations of zero!
  - computer must have two sets of instructions
    - one for signed arithmetic, one for unsigned



# Signed Integers

- **Ones' complement**

- negate a number by forming its bit-wise complement

- » e.g.,  $(-1) \cdot 01101011 = 10010100$

$b_{w-1} = 0 \Rightarrow$  non-negative number

$$\text{value} = \sum_{i=0}^{w-2} b_i \cdot 2^i$$

$b_{w-1} = 1 \Rightarrow$  negative number

$$\text{value} = \sum_{i=0}^{w-2} (b_i - 1) \cdot 2^i$$

two zeros!

# Signed Integers

- **Two's complement**

$b_{w-1} = 0 \Rightarrow$  non-negative number

$$\text{value} = \sum_{i=0}^{w-2} b_i \cdot 2^i$$

$b_{w-1} = 1 \Rightarrow$  negative number

$$\text{value} = (-1) \cdot 2^{w-1} + \sum_{i=0}^{w-2} b_i \cdot 2^i$$

one zero!

# Example

- $w = 4$

0000: 0

0001: 1

0010: 2

0011: 3

0100: 4

0101: 5

0110: 6

0111: 7

1000: -8

1001: -7

1010: -6

1011: -5

1100: -4

1101: -3

1110: -2

1111: -1

# Signed Integers

- **Negating two's complement**

$$value = -b_{w-1} 2^{w-1} + \sum_{i=0}^{w-2} b_i 2^i$$

- **how to compute  $-value$ ?**  
 **$(\sim value)+1$**

# Signed Integers

- Negating two's complement (continued)

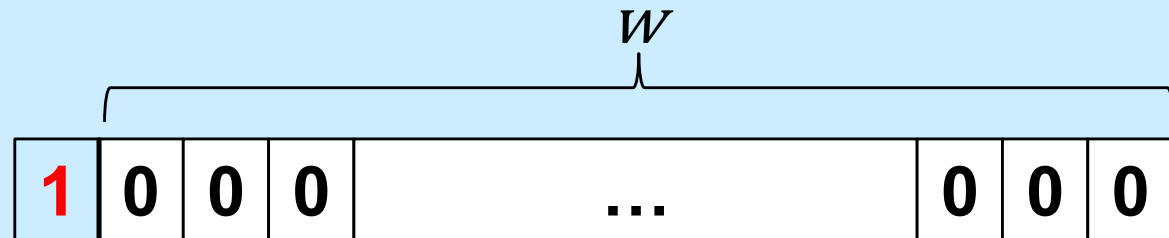
$$value + (\sim value + 1)$$

$$= (value + \sim value) + 1$$

$$= (2^w - 1) + 1$$

$$= 2^w$$

=



# Quiz 3

- **We have a computer with 4-bit words that uses two's complement to represent signed integers. What is the result of subtracting 0010 (2) from 0001 (1)?**
  - a) 1110
  - b) 1001
  - c) 0111
  - d) 1111

# Signed vs. Unsigned in C

- **char, short, int, and long**
  - signed integer types
  - right shift (>>) is arithmetic
- **unsigned char, unsigned short, unsigned int, unsigned long**
  - unsigned integer types
  - right shift (>>) is logical

# Numeric Ranges

- **Unsigned Values**

- $UMin = 0$

- $000\dots0$

- $UMax = 2^w - 1$

- $111\dots1$

- **Two's Complement Values**

- $TMin = -2^{w-1}$

- $100\dots0$

- $TMax = 2^{w-1} - 1$

- $011\dots1$

- **Other Values**

- Minus 1

- $111\dots1$

## Values for $W = 16$

	Decimal	Hex	Binary
<b>UMax</b>	<b>65535</b>	<b>FF FF</b>	<b>11111111 11111111</b>
<b>TMax</b>	<b>32767</b>	<b>7F FF</b>	<b>01111111 11111111</b>
<b>TMin</b>	<b>-32768</b>	<b>80 00</b>	<b>10000000 00000000</b>
<b>-1</b>	<b>-1</b>	<b>FF FF</b>	<b>11111111 11111111</b>
<b>0</b>	<b>0</b>	<b>00 00</b>	<b>00000000 00000000</b>



# Values for Different Word Sizes

	W			
	8	16	32	64
UMax	255	65,535	4,294,967,295	18,446,744,073,709,551,615
TMax	127	32,767	2,147,483,647	9,223,372,036,854,775,807
TMin	-128	-32,768	-2,147,483,648	-9,223,372,036,854,775,808

- **Observations**

$$|TMin| = TMax + 1$$

» Asymmetric range

$$UMax = 2 * TMax + 1$$

- **C Programming**

- `#include <limits.h>`
- declares constants, e.g.,
  - `ULONG_MAX`
  - `LONG_MAX`
  - `LONG_MIN`
- values platform-specific

# Quiz 4

- **What is  $-TMin$  (assuming two's complement signed integers)?**
  - a) **TMin**
  - b) **TMax**
  - c) **0**
  - d) **1**